

# Sorting Networks

A Lecture in CE Freshman Seminar Series:  
Ten Puzzling Problems in Computer Engineering



May 2020



Sorting Networks



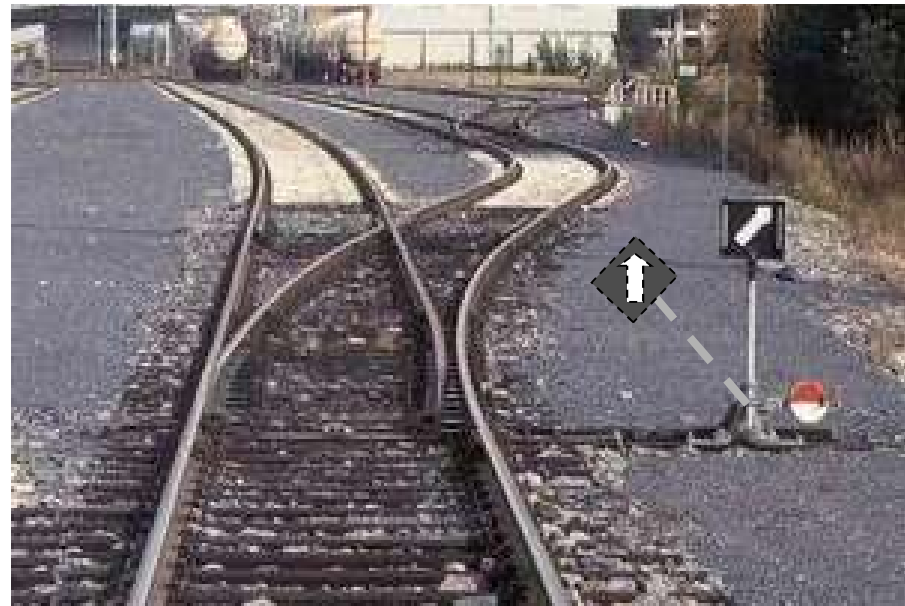
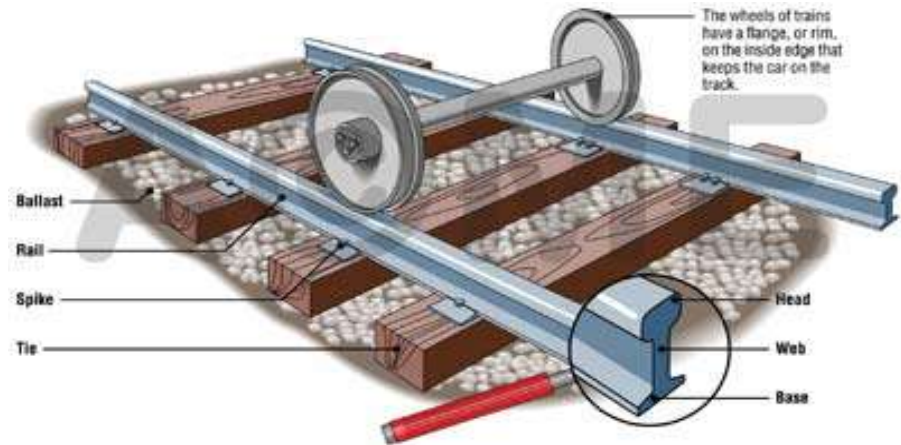
Slide 1

# About This Presentation

This presentation belongs to the lecture series entitled “Ten Puzzling Problems in Computer Engineering,” devised for a ten-week, one-unit, freshman seminar course by Behrooz Parhami, Professor of Computer Engineering at University of California, Santa Barbara. The material can be used freely in teaching and other educational settings. Unauthorized uses, including any use for financial gain, are prohibited. © Behrooz Parhami

| <b>Edition</b> | <b>Released</b> | <b>Revised</b>  | <b>Revised</b>  | <b>Revised</b>  | <b>Revised</b>  |
|----------------|-----------------|-----------------|-----------------|-----------------|-----------------|
| <b>First</b>   | <b>May 2007</b> | <b>May 2008</b> | <b>May 2009</b> | <b>May 2010</b> | <b>May 2011</b> |
|                |                 | <b>May 2012</b> | <b>May 2015</b> | <b>May 2016</b> | <b>May 2020</b> |
|                |                 |                 |                 |                 |                 |

# Railroad Tracks and Switches



# Coupling and Decoupling of Train Cars



Train cars and engines can be coupled and decoupled quickly



An engine can push a string of cars, or pull a desired subset by decoupling them from the rest

# Railroad Yards Have Many Tracks and Switches



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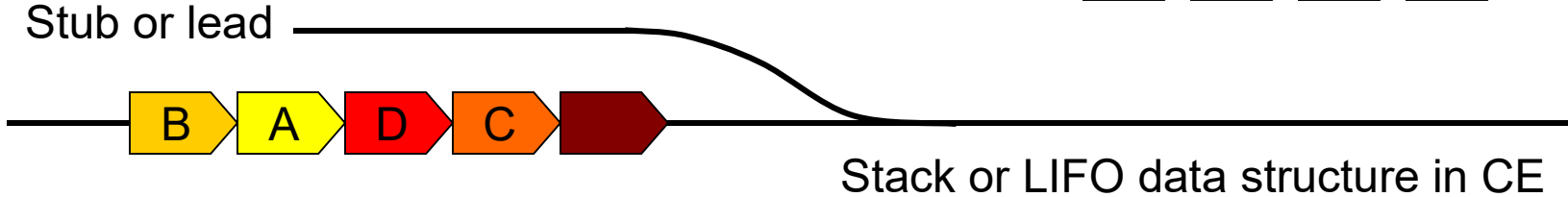
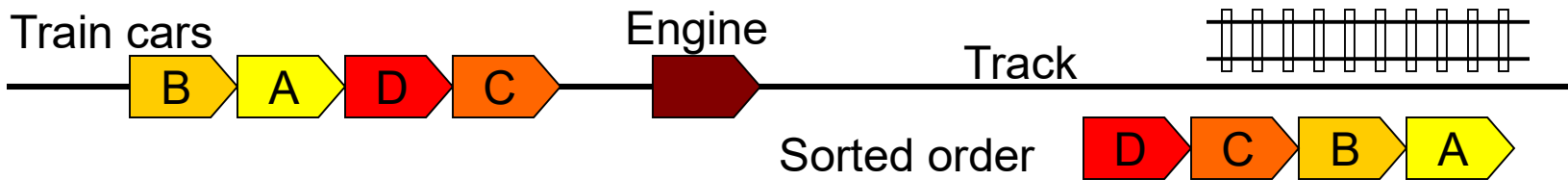


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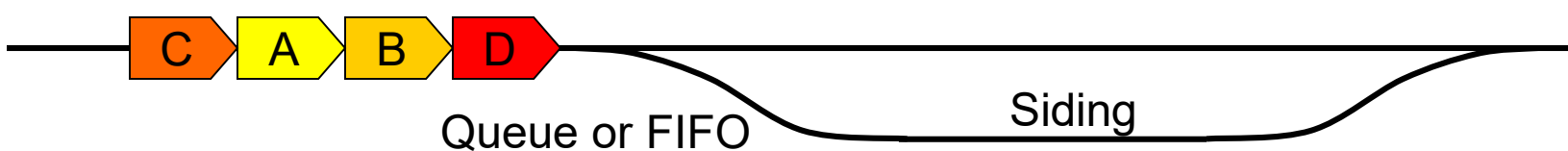
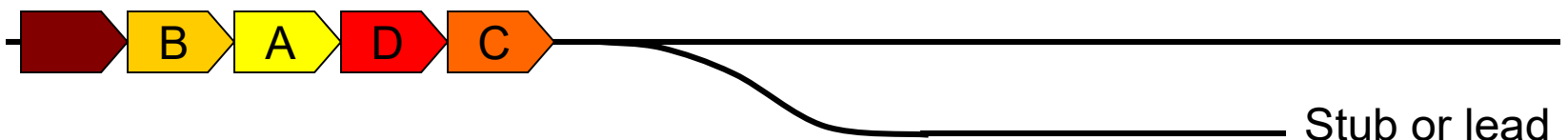
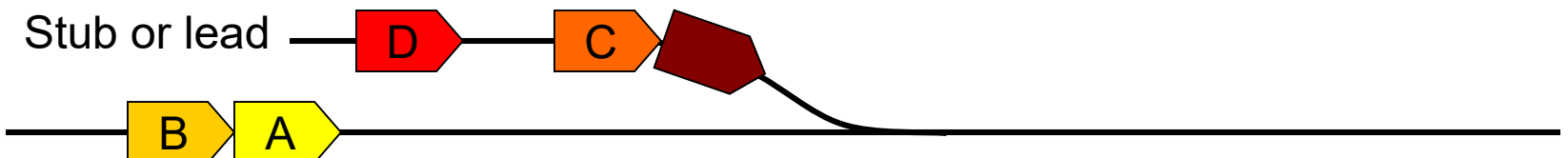


Slide 5

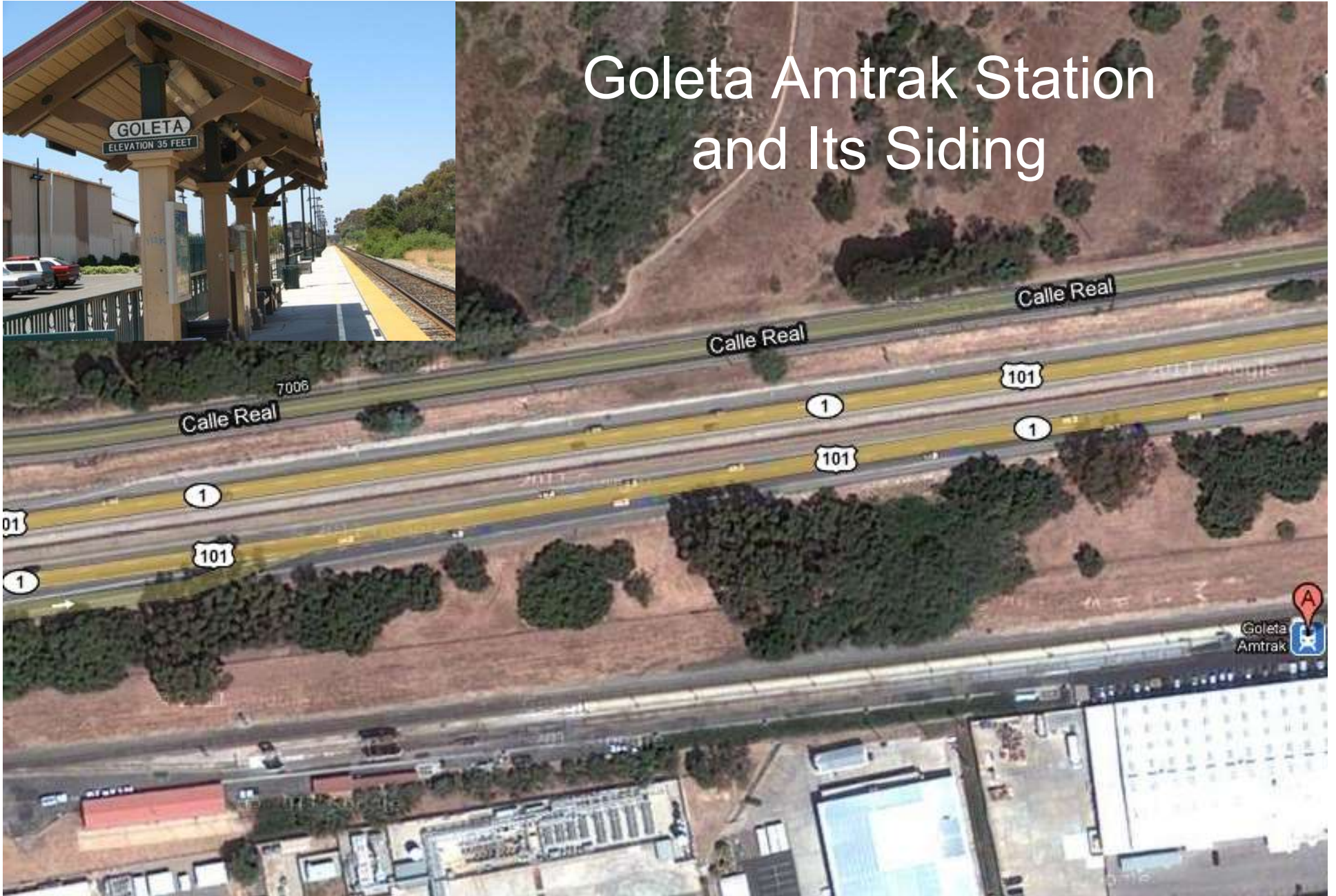
# Rearranging Trains



**Sorting algorithm:** Assemble train in stub, beginning with the last car  
 Repeat: If the next car is X, decouple train after X, push X into stub



# Goleta Amtrak Station and Its Siding



May 2020



Sorting Networks



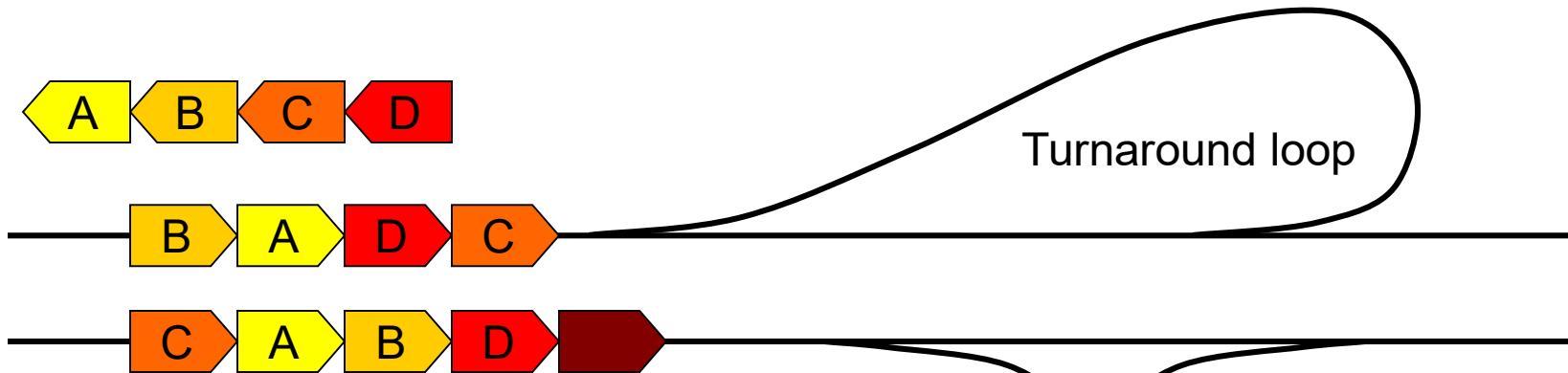
Slide 7

# Model Railroad Yard

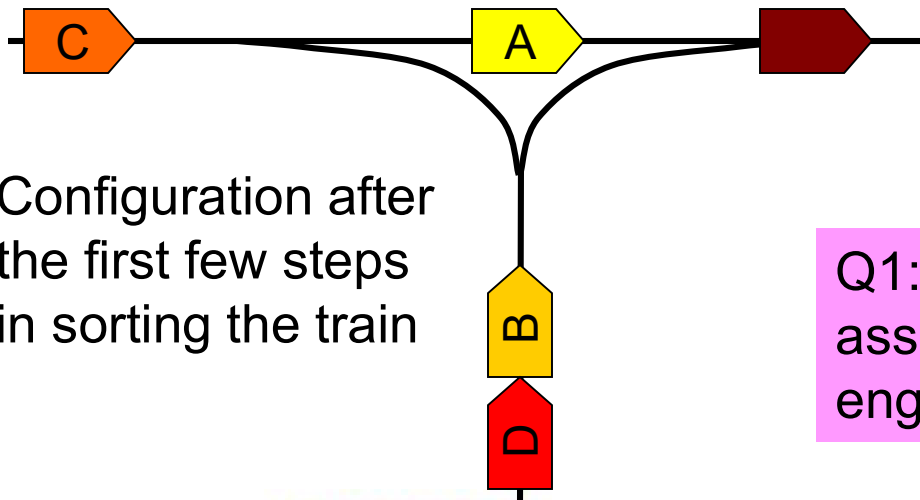




# Rearrangement with Change of Direction



What types of sorting are possible with a turnaround loop? A wye?



Configuration after the first few steps in sorting the train

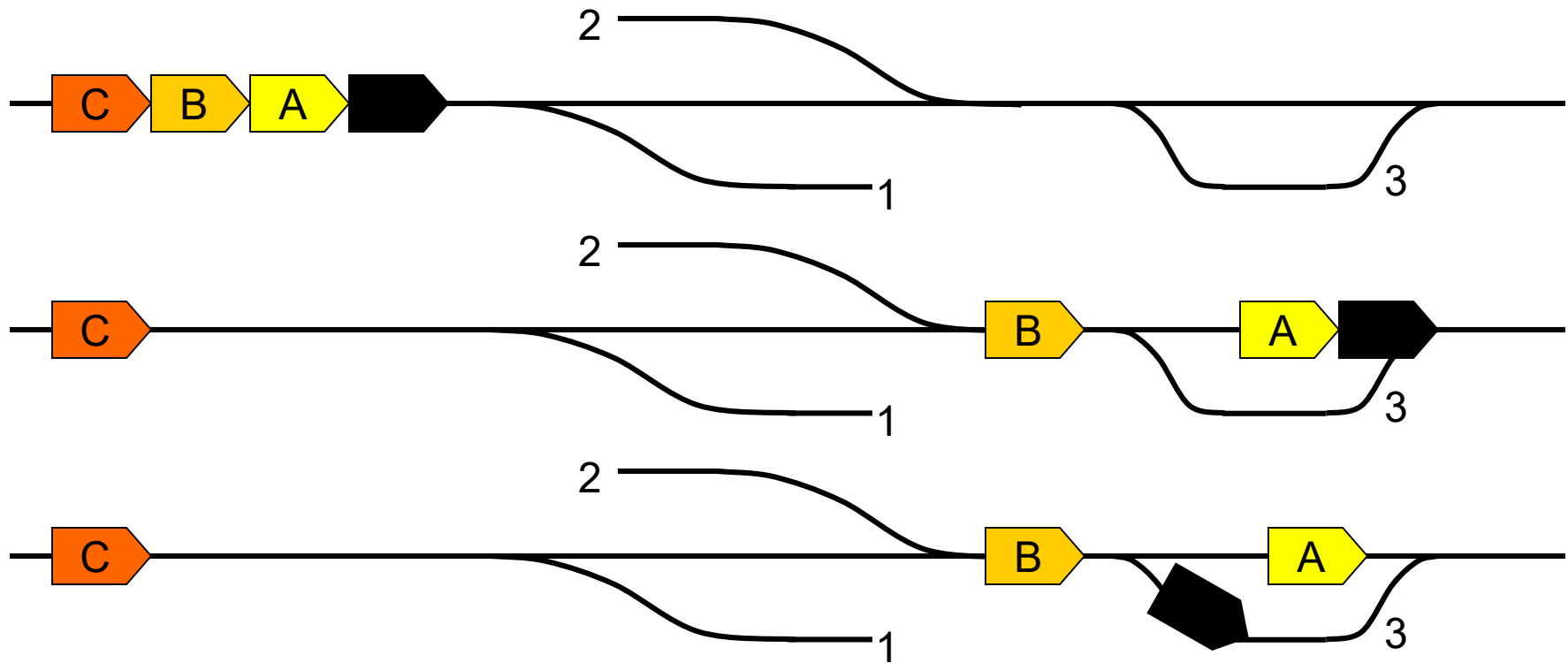
Q1: Complete the sorting process, assuming that it is okay for the engine to end up behind the train

# Model Train Turntable



# Delivering Train Cars in a Specific Order

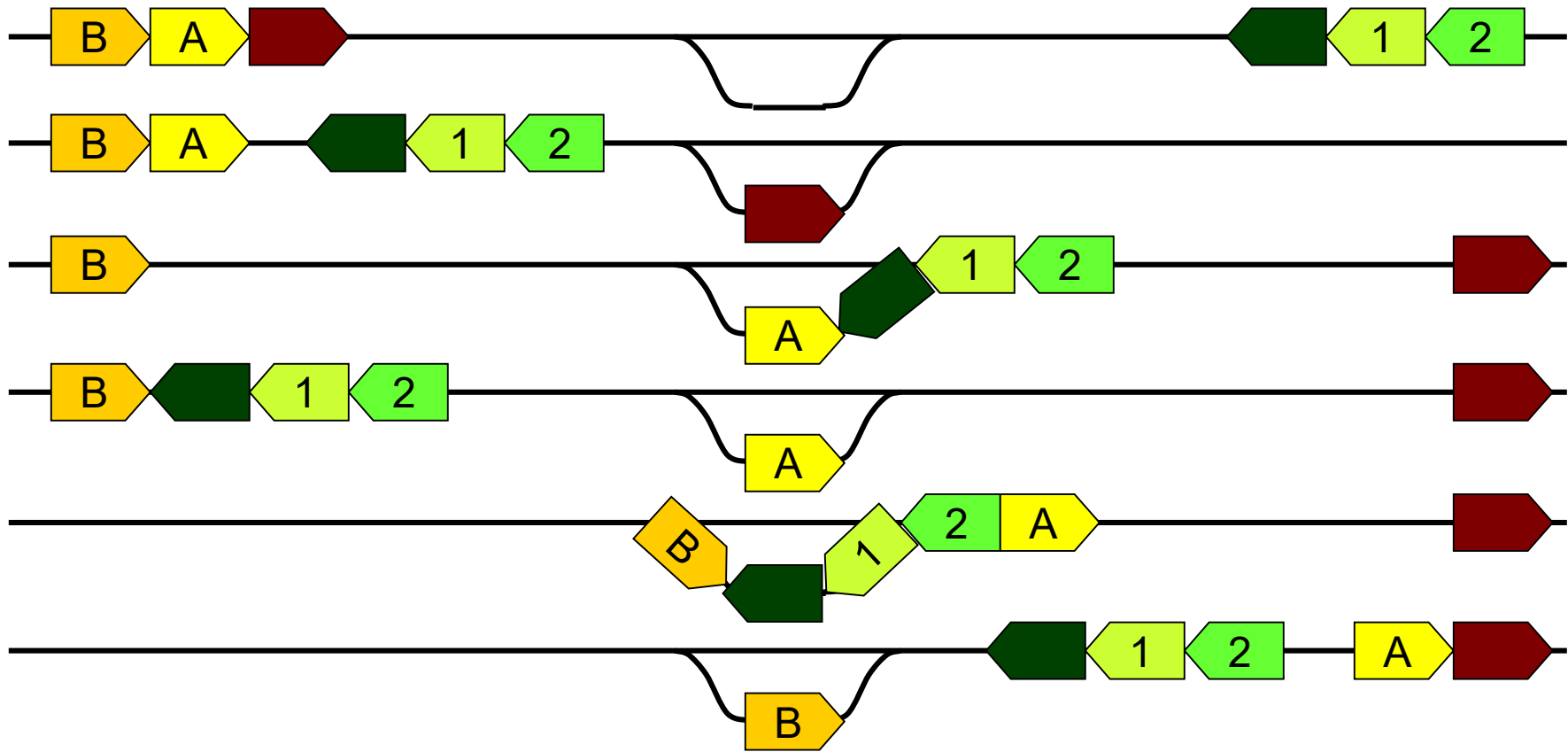
Cars in the train below have been sorted according to their delivery points. However, it is still nontrivial to deposit car A in stub 1, car B in stub 2, and car C in siding 3. Cars can be pulled or pushed by the engine.



Is there a better initial ordering of the cars for the deliveries in this puzzle?

# Train Passing Puzzle

The trains below must pass each other using a siding that can hold only one car or one engine. Show how this can be done.

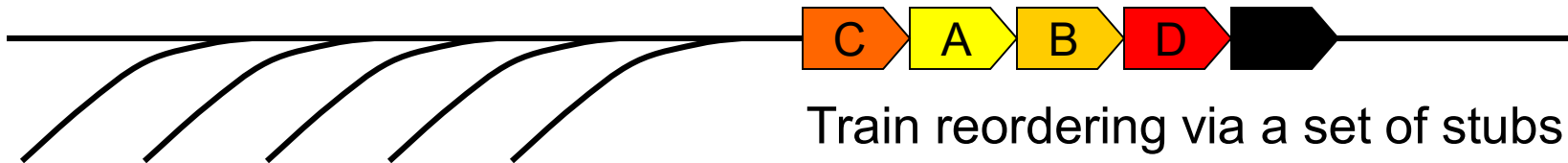


Q2: If the left and right trains have  $L$  and  $R$  cars, respectively, how many times will the siding be used for the trains to pass?

# Fast Combining or Reordering of Train Cars



Forming multiple trains  
from incoming cars



Train reordering via a set of stubs

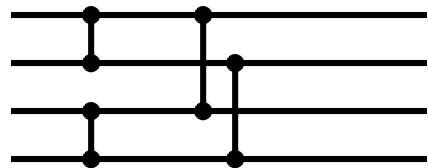
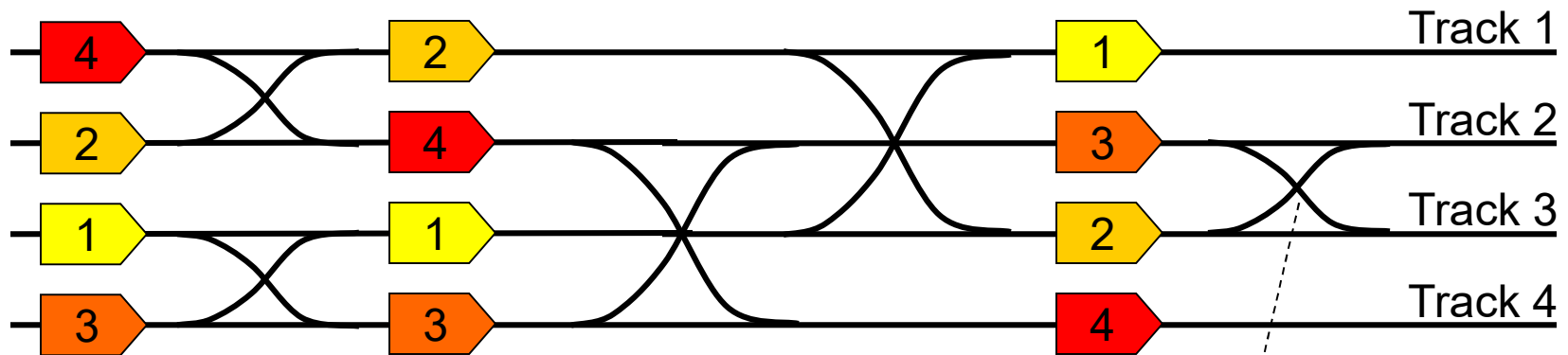
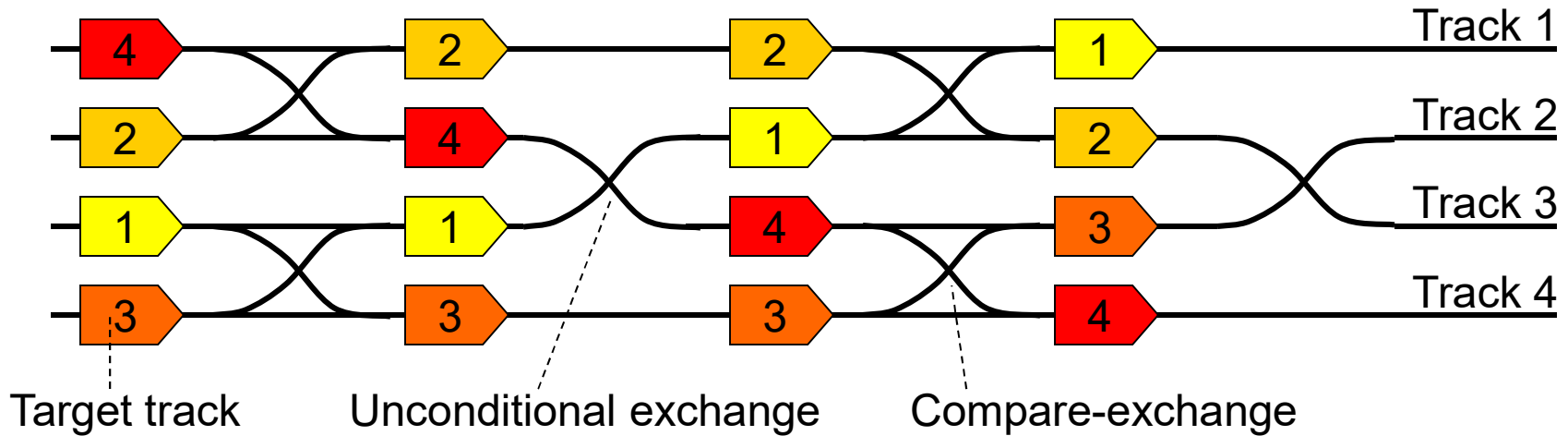


Cars are pushed or pulled by an engine  
Alternatively, movement in one direction  
may be achieved via sloped rails

Switches used to be adjusted manually,  
but nowadays, electronic control is used

Q3: Sort the two trains shown above  
with the three sidings and five stubs

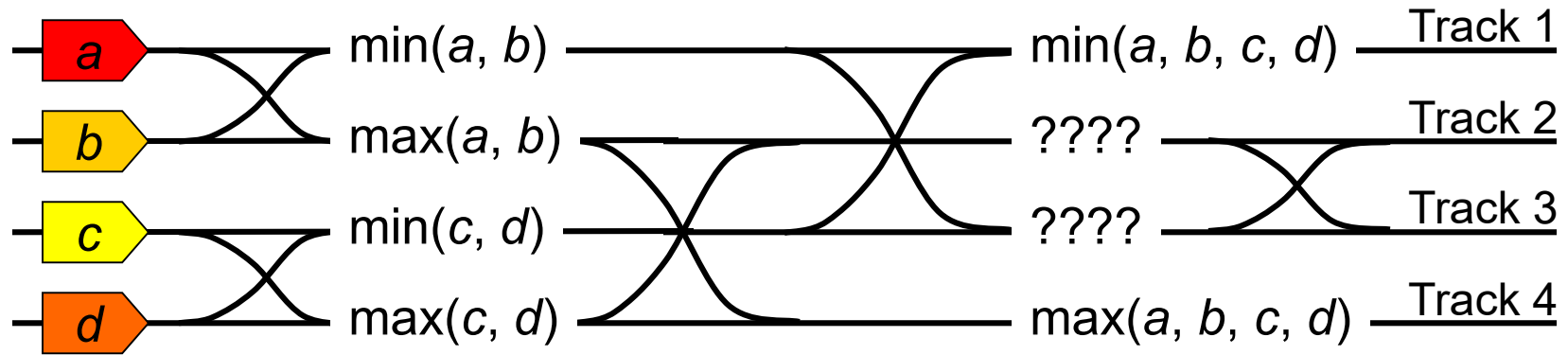
# Sorting Train Cars in Parallel



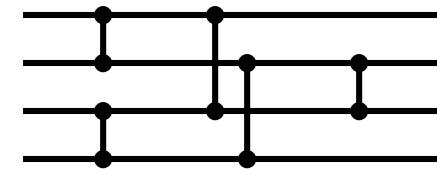
The net in stick diagram schematic

Is adding this compare-exchange element sufficient for producing a valid sorting network?

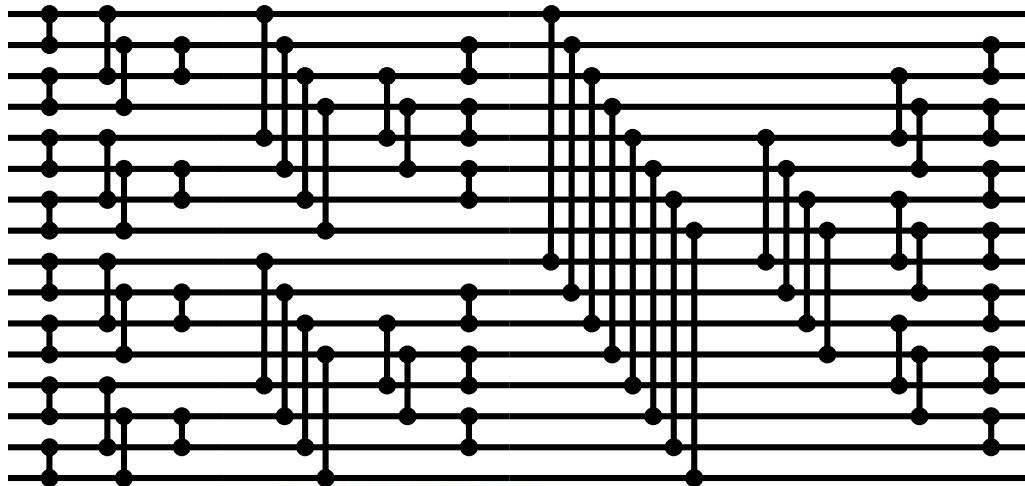
# Validating a Sorting Network



In the example above, it was fairly easy to show the validity of the sorting network. Generally, it is much more difficult



Stick diagram schematic

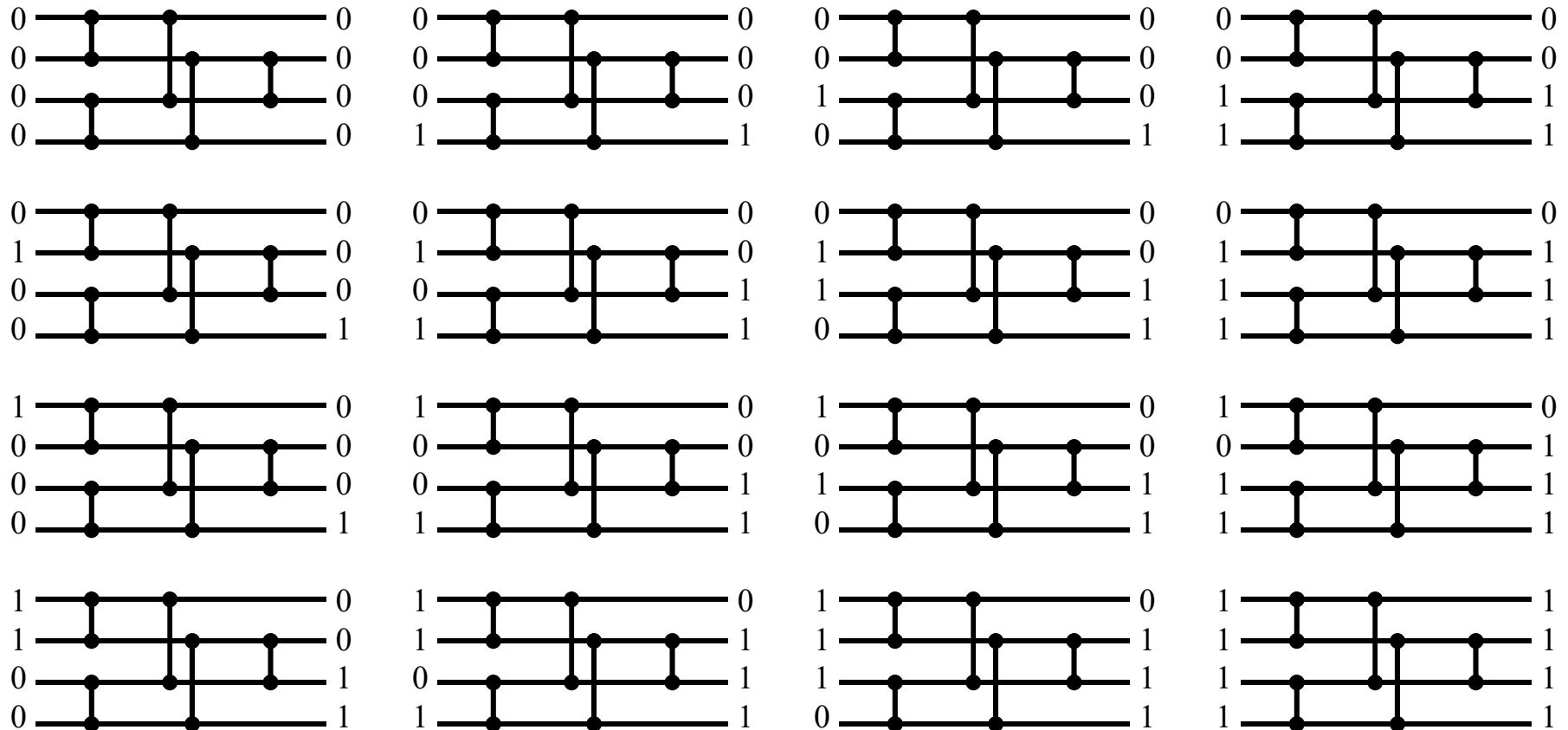


How would one establish the validity of this 16-input sorting network?

More importantly, how does one come up with this design in the first place?

# The Zero-One Principle

A sorting net built of comparators is valid if it correctly sorts all 0-1 sequences



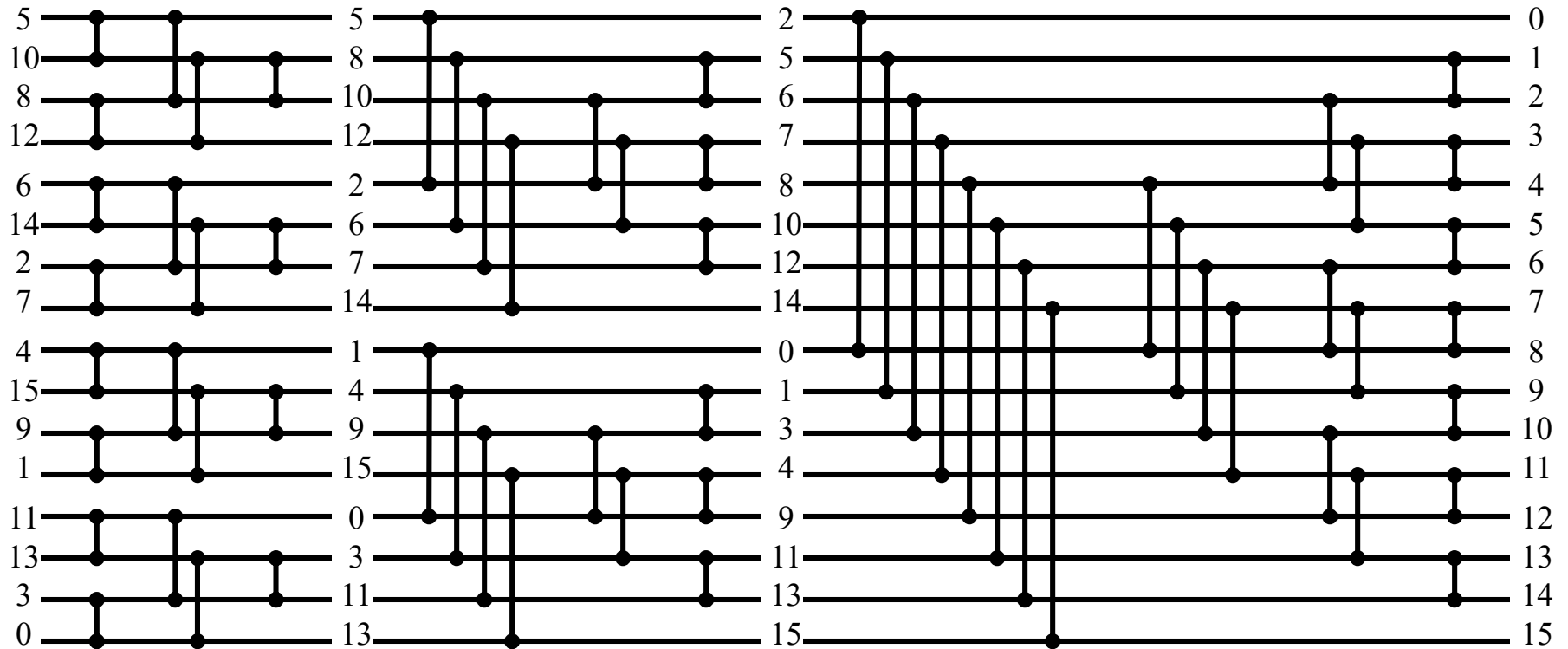
So, we can validate a sorting network using  $2^n$  rather than  $n!$  input patterns

$n = 12$ :  $2^n = 4096$ ,  $n! = 479,001,600$  (thousands vs. half a billion)



# A 16-Input Sorting Network

Use 4-input sorters, follow by (4, 4)-mergers, and end with an (8, 8)-merger



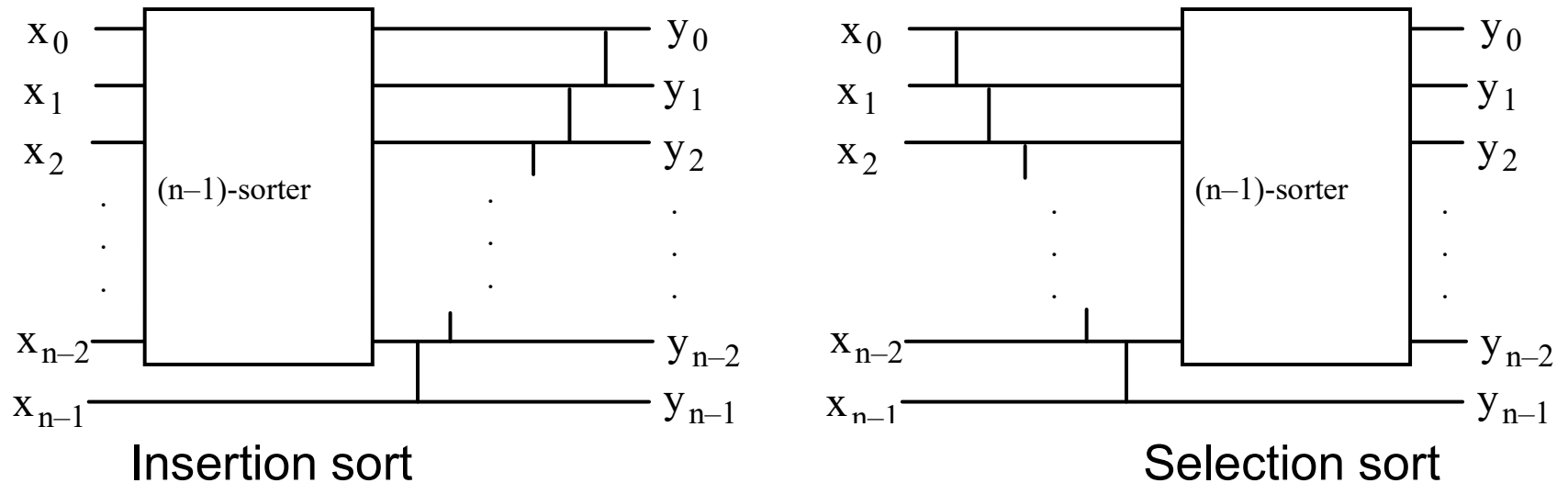
Using the 0-1 principle, we can validate this network via  $16 + 25 + 81$  tests

4-sorter tests

(4, 4)-merger tests

(8, 8)-merger tests

# Insertion Sort and Selection Sort



$$C(n) = n(n-1)/2$$

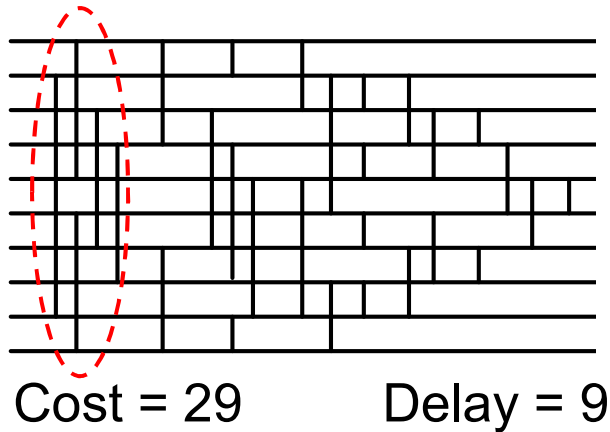
$$D(n) = 2n - 3$$

$$\text{Cost} \times \text{Delay} = \Theta(n^3)$$

|   |   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|---|
| 3 | 3 | 3 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 4 | 0 | 3 | 3 | 3 | 2 | 2 | 1 | 1 |
| 0 | 0 | 4 | 4 | 4 | 2 | 3 | 1 | 2 | 2 |
| 5 |   | 5 | 5 | 2 | 4 | 1 | 3 |   | 3 |
| 2 |   |   | 2 | 5 | 1 | 4 |   |   | 4 |
| 1 |   |   |   | 1 | 5 |   |   |   | 5 |

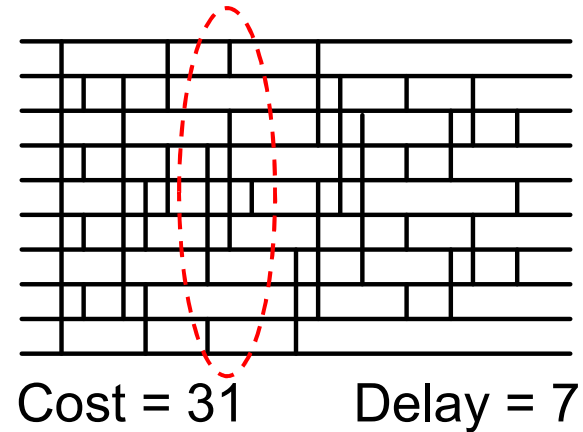
Fig. 7.8 Sorting network based on insertion sort or selection sort.

# The Best Sorting Networks



$$\text{Cost} \times \text{Delay} = 29 \times 9 = 261$$

Which  
10-input  
sorting  
network  
is better?



$$\text{Cost} \times \text{Delay} = 31 \times 7 = 217$$

Criterion 1: The number of sticks or compare-exchange blocks (cost)

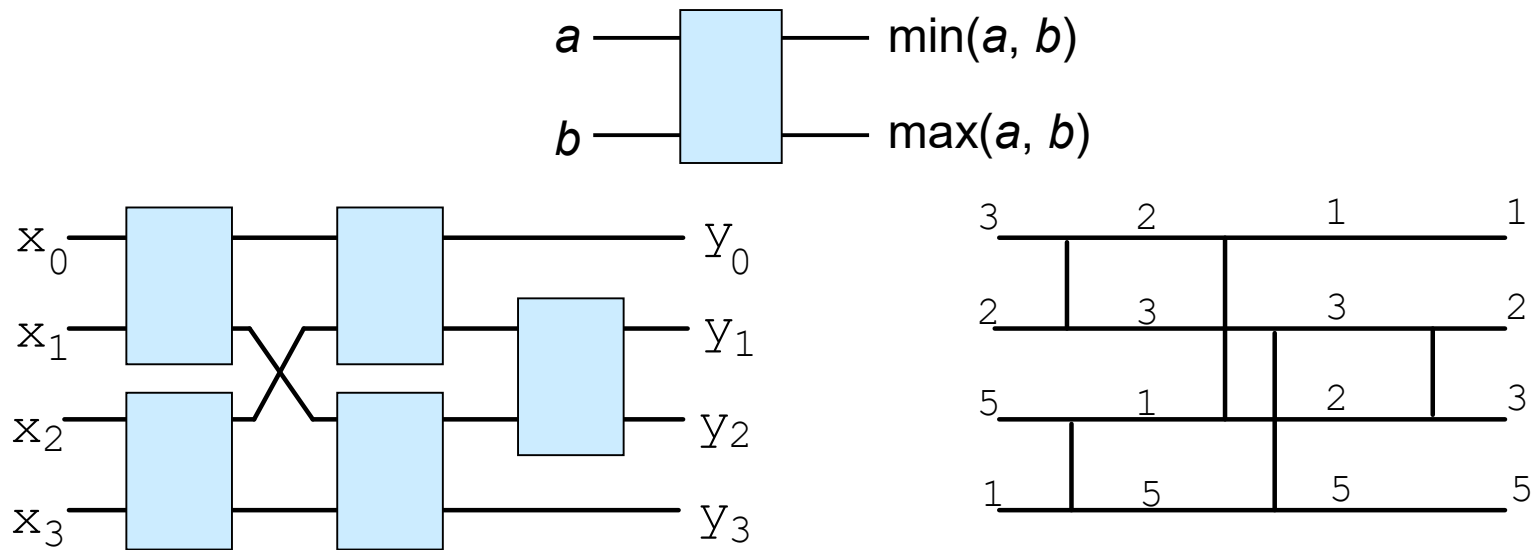
Criterion 2: The number of compare-exchanges in sequence (delay)

Criterion 3: The product of cost and delay (cost-effectiveness)

The most cost-effective  $n$ -input sorting network may be neither the fastest design, nor the lowest-cost design

# Electronic Sorting Networks

Electronic sorting networks are built of 2-sorters building blocks



## Applications of sorting networks:

- Directing information packets to their destinations in a network router
- Connecting  $n$  processors to  $n$  memory modules in a parallel computer

Q4: In the stick diagram of a 4-sorter on the top right, show that removing the top or bottom line and its comparators yields a valid 3-sorter but that removing one of the two middle lines does not