# Sequential Circuit Design with Verilog 

 ECE 152A - Fall 2006
## Reading Assignment

- Brown and Vranesic
- 6 Combinational - Circuit Building Blocks
- 6.6 Verilog for Combinational Circuits
- 6.6.1 The Conditional Operator
- 6.6.2 The If-Else Statement
- 6.6.3 The Case Statement


## Reading Assignment

- Brown and Vranesic (cont)
- 7 Flip-Flops, Registers, Counters, and a Simple Processor
- 7.12 Using Storage Elements with CAD Tools
- 7.12.2 Using Verilog Constructs for Storage Elements
- 7.12.3 Blocking and Non-Blocking Assignments
- 7.12.4 Non-Blocking Assignments for Combinational Circuits
- 7.12.5 Flip-Flops with Clear Capability
- 7.13 Using Registers and Counters with CAD Tools
- 7.13.3 Using Verilog Constructs for Registers and Counters


## The Gated D Latch

## - Transparent on high phase of clock

```
module D_latch(D, Clk, Q);
    input D, Clk;
    output Q;
    reg Q;
    always @(D or Clk)
        if (Clk)
            Q = D;
endmodule
```


## The Gated D Latch

- The "if" construct
- When D or CLK change value:
- if $C L K=1$, set $\mathrm{Q}=\mathrm{D}$
- Since there is no else, assignment occurs only when CLK = 1
- Q follows D when CLK $=1$
- Q remains latched on CLK = 0
- "Always" construct triggered by change in value of D or CLK
- Either change can cause the output to change


## The Gated D Latch

- The "always" construct
- Responds to changes in the signals on the sensitivity list
- always @ (D or Clk)
- Example above is "level sensitive"
- When D or Clk changes value
- Make edge triggered by using Verilog keywords posedge and negedge
- i.e., always @ (posedge Clk)


## The Edge Triggered D Flip-Flop

- Positive edge triggered

```
module flipflop(D, Clock, Q);
    input D, Clock;
    output Q;
    reg Q;
    always @(posedge Clock)
        Q = D; // Q+ = D, characteristic function
```

endmodule

## The Edge Triggered D Flip-Flop

- $D$ is not included on sensitivity list since it cannot cause output ( Q ) to change
- No transparent phase with edge triggered flipflops
- Characteristic function used in assignment statement
- Defining next state $\left(\mathrm{Q}^{+}\right)$of the flip-flop


## The Edge Triggered JK Flip-Flop

- Assign characteristic function to Q on rising clock edge ( $\mathrm{Q}^{+}=\mathrm{JQ}{ }^{\prime}+\mathrm{K}^{\prime} \mathrm{Q}$ )
module JKflipflop(J,K, Clock, Q);
input J,K, Clock;
output Q;
reg Q;
always @(posedge Clock)

$$
\mathrm{Q}=\mathrm{J} \& \& \sim \mathrm{Q} \| \sim \mathrm{K} \& \& \mathrm{Q} ; / / \mathrm{Q}^{+}=\mathrm{JQ} \mathrm{Q}^{\prime}+\mathrm{K}^{\prime} \mathrm{Q}
$$

endmodule

## The Edge Triggered JK Flip-Flop

- Functional Simulation



## The Edge Triggered T Flip-Flop

- Assign characteristic function to Q on rising clock edge ( $\mathrm{Q}^{+}=\mathrm{T}$ XOR Q)
module Tflipflop(T, Clock, Q);
input T, Clock;
output Q;
reg Q;
always @(posedge Clock)

$$
\mathrm{Q}=\mathrm{T} \wedge \mathrm{Q} ; \quad / / \mathrm{Q}=\mathrm{T} \text { XOR } \mathrm{Q}
$$

endmodule

## The Edge Triggered T Flip-Flop

- Functional Simulation



## Blocking and Non-Blocking Assignments

- Q = D
- Equal sign (=) signifies a blocking assignment
- Statements are evaluated in the order in which they are written
- If a variable is given a value by a blocking assignment, the new value is used in evaluating all subsequent statements in the block


## Blocking and Non-Blocking Assignments

- Blocking Assignment Statement Example

```
module example1(D, Clock, Q1, Q2);
    input D, Clock;
    output Q1, Q2;
    reg Q1, Q2;
    always @(posedge Clock)
    begin
        \(\mathrm{Q} 1=\mathrm{D}\);
        \(\mathrm{Q} 2=\mathrm{Q} 1\);
    end
endmodule
```


## Blocking and Non-Blocking Assignments

- Example synthesizes two positive edge triggered D flip-flops
- Both flip-flops triggered by same clock edge
- Both assignments in always block are blocking
- Q1 gets the value D
- Q2 then gets the new value of Q1
- Q1 ${ }^{+}$, which is now D

Blocking and Non-Blocking Assignments

- The synthesized circuit with blocking assignment statements



## Blocking and Non-Blocking Assignments

- Non-Blocking Statements (<=)
- Non-blocking assignment statements in an always block are evaluated using the values of the variables when the block is entered

```
always @(posedge Clock)
    begin
        Q1 <= D; // substitute
        Q2 <= Q1; // non-blocking assignments
    end
```

Blocking and Non-Blocking Assignments

- Q2 gets the value of Q1 when the always block is entered
- The synthesized circuit with non-blocking assignment statements
- Flip-flops connected in cascade



## Blocking and Non-Blocking Assignments

- Blocking Assignment Statement Example

```
module example3(x1, x2, x3, Clock, f, g);
    input x1, x2, x3, Clock;
    output f, g;
    reg f, g;
    always @(posedge Clock)
    begin
        f = x1 & x2;
        g = f | x3;
    end
endmodule
```

Blocking and Non-Blocking Assignments

- Both fand $g$ are implemented as the outputs of D flip-flops
- Synthesized as flip-flops because the sensitivity list of the always block specifies posedge Clock
- " $g$ " gets the new value $\left(\mathrm{Q}^{+}\right)$of " f " OR'd with $\mathrm{x}_{3}$


## Blocking and Non-Blocking Assignments

- The synthesized circuit
- Blocking assignment statements


Blocking and Non-Blocking Assignments

- If assignment statements changed to nonblocking

```
always @(posedge Clock)
    begin
            f <= x1 & x2;
            g <= f | x 3;
    end
```

- " $g$ " gets the previous value of " $f$ " (the value when the always block is entered, i.e., Q)


## Blocking and Non-Blocking Assignments

## - The synthesized circuit

- Non-blocking assignment statements


Blocking and Non-Blocking Assignments

- General Rule
- The results of non-blocking assignments are visible only after all of the statements in the always block have been evaluated
- When there are multiple assignments to the same variable inside an always block, the result of the last assignment is maintained


## Flip-Flops with Clear

## - Asynchronous Clear

```
module flipflop(D, Clock, Resetn, Q);
    input D, Clock, Resetn;
    output Q;
    reg Q;
    always @(negedge Resetn or posedge Clock)
        if (!Resetn)
        Q<= 0;
        else
            Q <= D;
```

endmodule

## Flip-Flops with Clear

- Synchronous Clear
module flipflop(D, Clock, Resetn, Q);
input D, Clock, Resetn;
output Q;
reg Q;
always @(posedge Clock)
if (!Resetn) // check value of reset on clock edge
$\mathrm{Q}<=0$;
else
$\mathrm{Q}<=\mathrm{D}$;
endmodule


## 4-Bit Binary Counter

- Counter includes reset and enable

```
module upcount(Resetn, Clock, E, Q);
    input Resetn, Clock, E;
        output [3:0] Q;
    reg [3:0] Q;
    always @(negedge Resetn or posedge Clock)
        if (!Resetn)
            Q <= 0; // asynchronous reset overrides enable
        else if (E)
            \(\mathrm{Q}<=\mathrm{Q}+1\); // synthesizes adder circuit
```

endmodule

## 4-Bit Binary Counter

- Functional Simulation



## Finite State Machine (FSM) Design

- Recall state diagram for JK flip-flop counter from previous lecture


Finite State Machine (FSM) Design

- The State Table

|  | PS |  |  | NS |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A | B | C | $\mathrm{A}^{+}$ | $\mathrm{B}^{+}$ | $\mathrm{C}^{+}$ |
| 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | X | X | X |
| 0 | 1 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | X | X | X |
| 1 | 1 | 0 | X | X | X |
| 1 | 1 | 1 | 0 | 1 | 0 |

## Finite State Machine (FSM) Design

- Can't use addition operator because sequence is not binary count
- See previous example
- Use parameter statement to define states

$$
\begin{aligned}
\text { parameter }[2: 0] \mathrm{A} & =3^{\prime} \mathrm{b} 000, \mathrm{~B}=3^{\prime} \mathrm{b} 100, \mathrm{C}=3^{\prime} \mathrm{b} 111, \\
\mathrm{D} & =3^{\prime} \mathrm{b} 010, \mathrm{E}=3^{\prime} \mathrm{b} 011 ;
\end{aligned}
$$

## Finite State Machine (FSM) Design

- Use case statement to implement state transitions

```
always @ (posedge clock)
case(count)
        A: count <= B;
        B: count <= C;
        C}\mathrm{ : count <= D;
        D: count <= E;
        E: count <= A;
        default: count <= A;
endcase
```


## Finite State Machine (FSM) Design

- The complete module

```
module jk_counter(count, clock);
input clock;
output [2:0] count;
reg [2:0] count;
parameter [2:0] A = 3'b000, B = 3'b100, C = 3'b111,
        D=3'b010,E = 3'b011
always @ (posedge clock)
case(count)
        A: count <= B;
        B: count <= C
        C: count <= D;
        D: count <= E
        E: count <= A;
        default: count <= A
endcase
endmodule
```


## Finite State Machine (FSM) Design

- Functional Simulation


