Wireless Communications Lecture 10

[SNR per symbol and SNR per bit]

$$SNR = \frac{P_R}{N_0 B} = \frac{E_s}{N_0 B T_s} = \frac{E_b}{N_0 B T_b}$$

For BPSK:

$$T_b = T_s, \quad E_b = E_s, \quad \text{and} \quad T_s = 1/B.$$

Raised cosine pulse shaper for other pulses. $T_s \sim \frac{1}{B}$ The SNR becomes

$$SNR = \frac{E_s}{N_0} \text{ or } (\frac{E_b}{N_0} \text{ BPSK})$$

The SNR per symbol γ_s and the SNR per bit γ_b are defined as

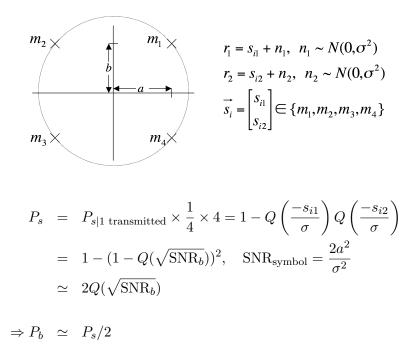
$$\gamma_s = \frac{E_s}{N_0}$$
$$\gamma_b = \frac{E_b}{N_0}$$

[Symbol error rate and bit error rate]

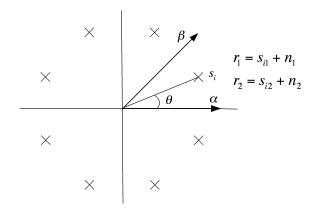
One approximation:

$$P_s \simeq$$
 one bit in error
= $P_{b_1}(1 - P_{b_2})(1 - P_{b_3}) \dots + P_{b_2}(1 - P_{b_1})(1 - P_{b_3}) \dots +$
= MP_{b_1} , M : the number of bits

² [QPSK]



[General case] For MPSK



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 $\implies \vec{r} = \vec{s}_i + \vec{n} \Rightarrow$ need angle distribution.

In-phase and quadrature Gaussian \Rightarrow bivariate distribution of |r| and $\theta \Rightarrow$ Integrate for θ distribution.

$$P_s = 1 - \int_{\alpha}^{\beta} P(\theta) d\theta$$

Closed form does not exist. The nearest neighbor approximation

$$P_s \simeq 2Q(\sqrt{2}\sqrt{\mathrm{SNR}_b})\sin(\pi/M)$$

[Average symbol error probability(general)] Many systems have : $P_s \simeq \alpha_M Q(\sqrt{\beta_M \gamma_s})$ Finding the average probability of symbol error

$$\begin{split} \bar{P}_s &= \int_0^\infty \alpha_M Q(\sqrt{\beta_M \gamma}) \lambda e^{-\lambda \gamma} d\gamma, \quad \lambda = \frac{1}{\bar{\gamma}_s} \\ &= \frac{\alpha_M}{2} \left[1 - \sqrt{\frac{0.5 \beta_M \bar{\gamma}_s}{1 + 0.5 \beta_M \bar{\gamma}_s}} \right] \\ &\approx \frac{\alpha_M}{2\beta_M \bar{\gamma}_s} \end{split}$$

Show P_b for AWGN and fading channel (figure 6.1, 6.2 Goldsmith)

[Alternative form of *Q*-function by Craig]

$$Q(z) = \frac{1}{\pi} \int_0^{\pi/2} e^{\frac{-z^2}{2\sin^2\varphi}} d\varphi, \quad z > 0$$

[Moment Generating Function]

$$\mathcal{M}_{\gamma}(s) = \int_{0}^{\infty} p_{\gamma}(\gamma) e^{s\gamma} d\gamma$$

Setting $\alpha = \alpha_M$ and $g = 0.5\beta_M$ in $P_s = \alpha_M Q(\sqrt{\beta_M \gamma_s})$

$$\begin{array}{lcl} P_s &=& \alpha_M Q(\sqrt{\beta_M \gamma_s})|_{\alpha = \alpha_M, g = 0.5 \beta_M} \\ P_s &=& \alpha Q(\sqrt{2g \gamma_s}) \\ && \text{Using the alternative Q-function} \\ &=& \frac{\alpha}{\pi} \int_0^{\pi/2} e^{\frac{-g \gamma_s}{\sin^2 \psi}} d\varphi \end{array}$$

The average error probability in fading for modulation with $P_s = \alpha Q(\sqrt{2g\gamma_s})$ is

$$\bar{P}_{s} = \frac{\alpha}{\pi} \int_{0}^{\infty} \int_{0}^{\pi/2} e^{\frac{-g\gamma}{\sin^{2}\varphi}} d\varphi p_{\gamma_{s}}(\gamma) d\gamma$$
$$= \frac{\alpha}{\pi} \int_{0}^{\pi/2} \left[\int_{0}^{\infty} e^{\frac{-g\gamma}{\sin^{2}\varphi}} p_{\gamma_{s}}(\gamma) d\gamma \right] d\varphi$$
$$= \frac{\alpha}{\pi} \int_{0}^{\pi/2} \mathcal{M}_{\gamma_{s}} \left(\frac{-g}{\sin^{2}\varphi} \right) d\varphi$$

The MGF $\mathcal{M}_{\gamma_s}(s)$ for Rayleigh distribution is

$$\mathcal{M}_{\gamma_s}(s) = (1 - s\bar{\gamma}_s)^{-1}, \quad s = \frac{-g}{\sin^2\varphi}$$

[Example] Find outage probability of BPSK at $P_b = 10^{-3}$ for Rayleigh fading with selection diversity for M = 2. Equal branch $\overline{\text{SNR}}_s$ of $\bar{\gamma} = 15$ dB.

$$P_b = 10^{-3} \Rightarrow \gamma_b = 10 \text{ dB} \Rightarrow \gamma_0 = 10$$

 $P_{\text{out}} = [1 - e^{-\gamma_0/\bar{\gamma}}]^2 = [1 - e^{-10/10^{1.5}}]^2$

[Average probability of symbol error(MRC)] The distribution of γ_{Σ} for MRC

$$p_{\gamma_{\Sigma}}(\gamma) = \frac{\gamma^{M-1} e^{-\gamma/\bar{\gamma}}}{\bar{\gamma}^M (M-1)!}, \quad \gamma \ge 0.$$

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The average probability of symbol error becomes

$$\bar{P}_s = \int_0^\infty \alpha Q(\sqrt{\beta\gamma}) p_{\gamma\Sigma}(\gamma) d\gamma$$

Capacity of fading channels

- Capacity of AWGN channel: $C = B \log_2(1 + \gamma)$ C: bits/sec, B:Bandwidth, $\gamma = \frac{P}{N_o B}$ (received SNR)
- Meaning of capacity
- Fading channels: SNR is a random variable.
- Finding capacity if only channel distribution is known @TX and RX is hard.
- Channel capacity with RX CSI (Channel Side Information)
 - RX knows the channel.
 - TX sends a fixed rate
 - Two types of capacity: Shannon and Outage
 - Shannon capacity: ergodic capacity, max rate with which you can still decode with zero error probability.
 - Capacity with outage: you allow some probability of error.

[Shannon (Ergodic) Capacity] For RX CSI, we will have

$$C = \int_0^\infty B \log_2(1+\gamma) p(r) dr$$

log is a concave function \Rightarrow applying Jenson's Inequality. For concave f(x),

$$\mathbf{E}(f(x)) \le f(\mathbf{E}(x)) \Rightarrow C = \int B \log_2(1+\gamma) p(\gamma) d\gamma \le B \log_2(1+\bar{\gamma})$$

It is worse than AWGN with $\bar{\gamma}$.

[Capacity with outage] RX CSI, outage case:

RX assumes γ_{min} for received SNR & sends

$$C = B \log_2(1 + \gamma_{min}) \Rightarrow \text{ if } \gamma < \gamma_{min} \Rightarrow \text{ bits cannot be decoded without error}$$

$$P_{out} = \text{Probability of not decoding correctly}$$

$$= P(\gamma < \gamma_{min})$$

$$C_{out} = \text{number of received correctly} = (1 - P_{out})B \log_2(1 + \gamma_{min})$$

 $\gamma_{min} \uparrow \rightarrow C \uparrow$ but

$$\gamma_{min} \uparrow \Rightarrow \left\{ \begin{array}{c} 1 + \gamma_{min} & \uparrow \\ 1 - P_{out} & \downarrow \end{array} \right.$$

Choose γ_{min} that gives the optimum.

Channel side information at Transmitter and Receiver

[Shannon capacity, TX & RX CSI:]

• Adapt the rate to SNR.

$$C = \max \int_0^\infty B \log_2 \left(1 + \frac{P(\gamma)\gamma}{\bar{P}} \right) p(\gamma) d\gamma$$
$$\int_0^\infty P(\gamma) p(\gamma) d\gamma \le \bar{P}$$

• Optimum $P(\gamma)$: Form Lagrangian :

$$J = \int_0^\infty B \log_2\left(1 + \frac{P(\gamma)\gamma}{\bar{P}}\right) p(\gamma)d\gamma - \lambda\left(\int_0^\infty P(\gamma)p(\gamma)d\gamma - \bar{P}\right)$$

differentiate with respect to $P(\gamma)$,

$$\frac{\partial J}{\partial P(\gamma)} = \frac{(B/\ln 2)(\gamma/\overline{P})}{1 + P(\gamma)\gamma/\overline{P}}p(\gamma) - \lambda p(\gamma) = 0$$

then,

$$P(\gamma) = \frac{(B/\ln 2)(\gamma/P) - \lambda}{\gamma(\lambda/\overline{P})} > 0$$
$$\gamma > \frac{\overline{P}\lambda}{B} \ln 2 = \gamma_0 \Rightarrow P(\gamma) = \begin{cases} \frac{\overline{P}}{\gamma_0} - \frac{\overline{P}}{\gamma} & \gamma \ge \gamma_0\\ 0 & \text{else} \end{cases}$$

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$$C = \int_{\gamma_0}^{\infty} B \log_2(\gamma/\gamma_0) p(\gamma) d\gamma$$

Finding γ_0 :

$$\int_{0}^{\infty} \frac{P(\gamma)}{\overline{P}} p(\gamma) d\gamma = 1$$
$$\int_{\gamma_{0}}^{\infty} \left(\frac{1}{\gamma_{0}} - \frac{1}{\gamma}\right) p(\gamma) d\gamma = 1$$

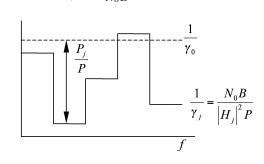
[Wideband Channels:] want to distribute total TX power over the frequency bins:

$$C = \sum B \log_2 \left(1 + \frac{|H_j|^2 P_j}{N_o B} \right)$$
$$\sum P_j \leq P$$

Similar Lagrangian method:

$$\frac{P_j}{P} = \begin{cases} \frac{1}{\gamma_0} - \frac{1}{\gamma_j} & \gamma_j \ge \gamma_0\\ 0 & \text{else} \end{cases}$$

Find $\gamma_0 \Rightarrow \sum_j \frac{1}{\gamma_0} - \frac{1}{\gamma_j} = 1$ where $\gamma_j = \frac{|H_j|^2 P_j}{N_o B}$



Water-filling (Multicarrier Modulation) : DSL